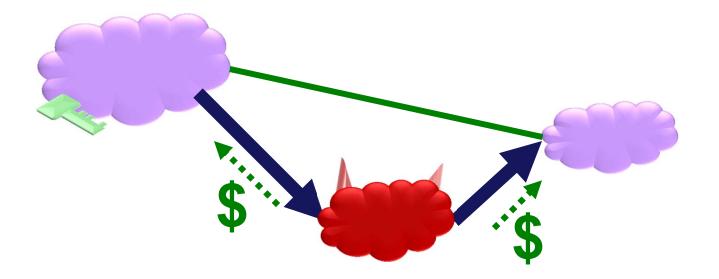
How Secure are

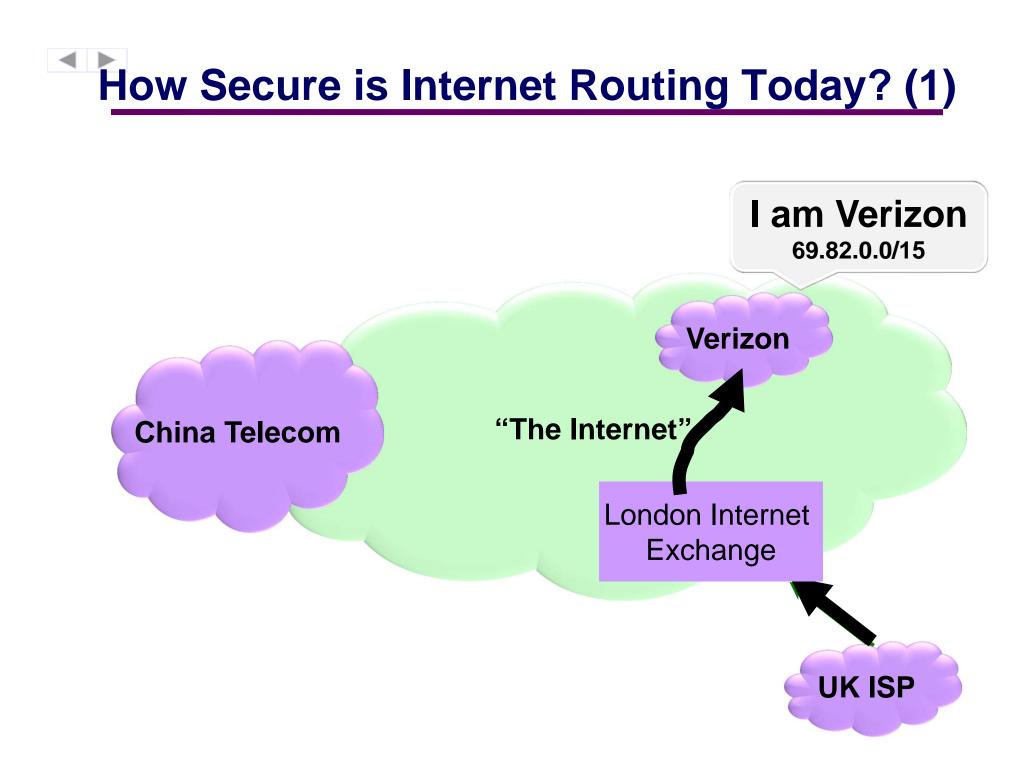
UCSD Systems Seminar February 10, 2011

Secure Internet Routing Protocols?



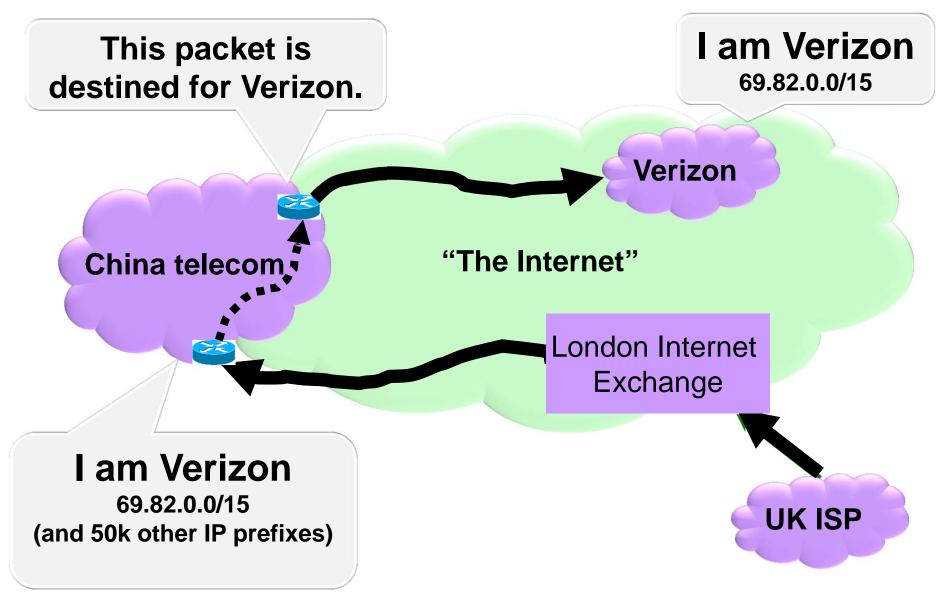
Sharon Goldberg Boston University

Michael Schapira Princeton Pete Hummon AT&T Jennifer Rexford Princeton



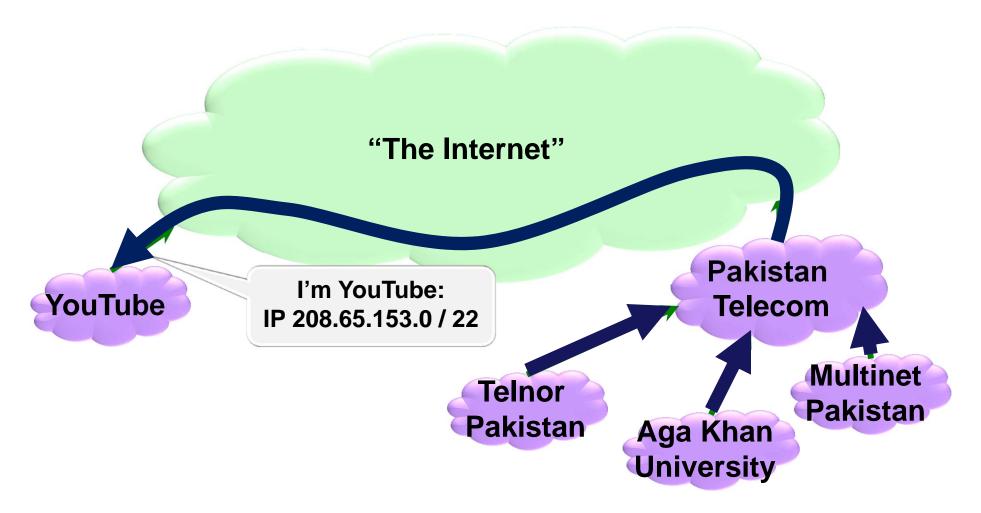
How Secure is Internet Routing Today? (2)

April 2010 : China Telecom intercepts traffic



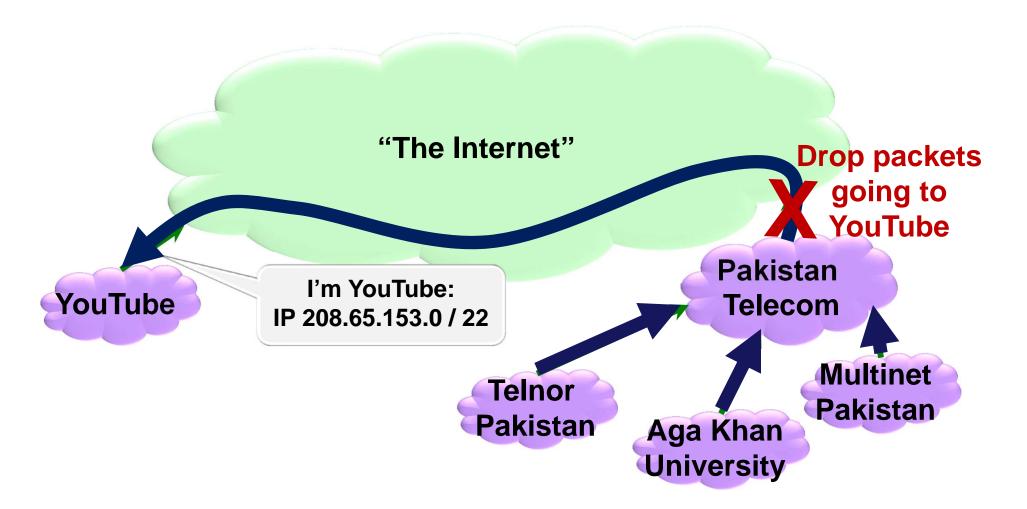
How Secure is Routing on the Internet Today? (3)

February 2008 : Pakistan Telecom hijacks Youtube



How Secure is Routing on the Internet Today? (4)

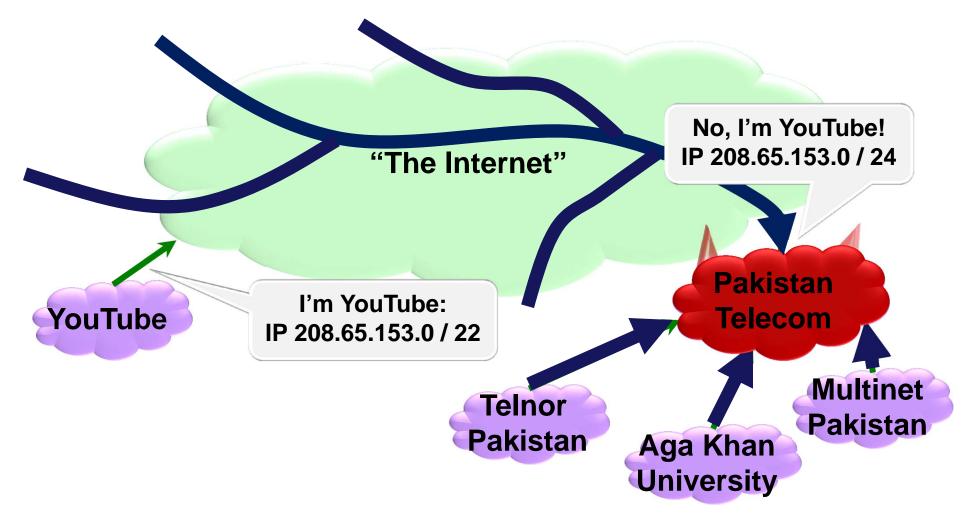
Here's what should have happened....



Block your own customers.

How Secure is Routing on the Internet Today? (5)

But here's what Pakistan ended up doing...



Draw traffic from the entire Internet!

Today, Internet routing is surprisingly insecure

- Decade of research on secure routing protocols
- With RPKI we can finally consider deploying one.

Our Goal: Compare the effectiveness of these protocols.

- Each has a different set of security properties.
- How well do they prevent attacks?

Our approach: Evaluate via simulation on network data.

- Data: Map of Internet & business relationships
- ... from [CAIDA] and [UCLA Cyclops]
- To compare protocols, we must find worst-case attacks









This talk

Pakistan Telecom hijacks YouTube

How Internet Routing Works

(and why economics matter)

Secure Routing Protocols and Attacks

Theory Interlude

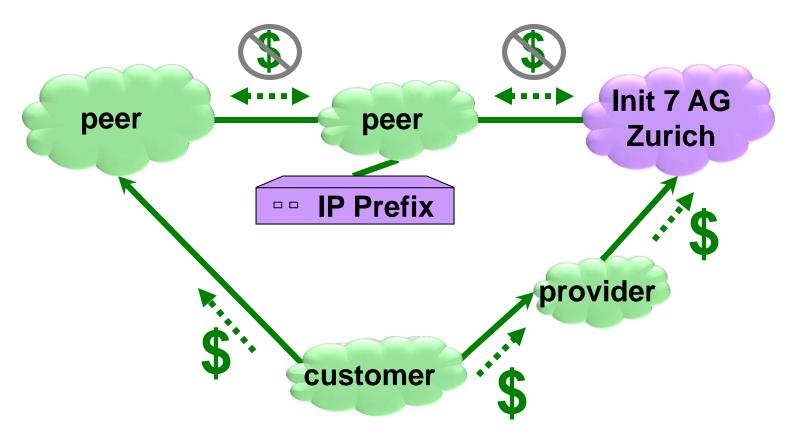
Results!



Implications & Deployment Challenges

BGP: The Internet's Routing Protocol (1)

The Border Gateway Protocol (BGP) sets up paths from Autonomous Systems (ASes) to destination IP addresses.

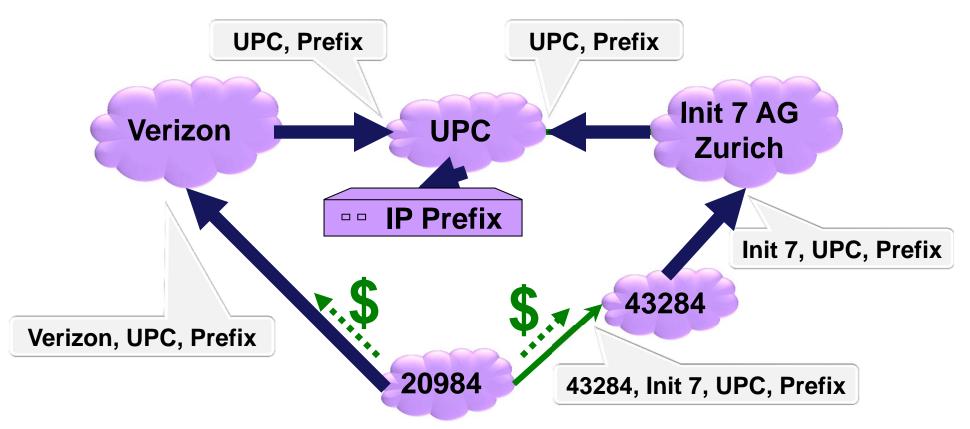


A model of routing decisions:

• Prefer cheaper paths. Then, prefer shorter paths.

BGP: The Internet's Routing Protocol (2)

The Border Gateway Protocol (BGP) sets up paths from Autonomous Systems (ASes) to destination IP addresses.

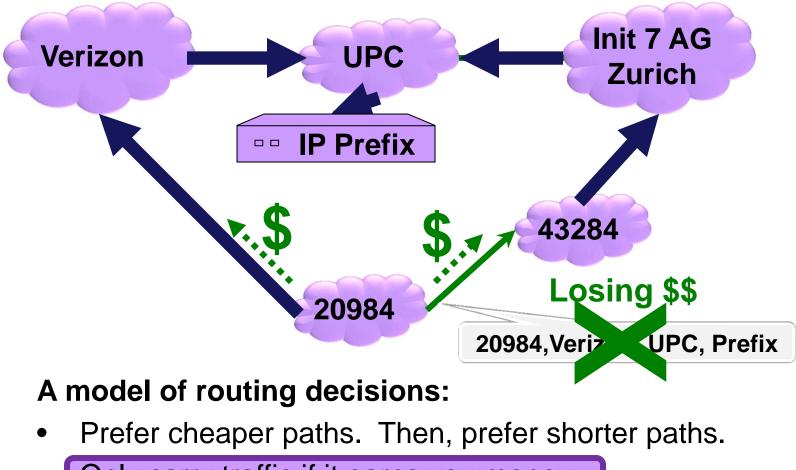


A model of routing decisions:

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BGP: The Internet's Routing Protocol (3)

The Border Gateway Protocol (BGP) sets up paths from Autonomous Systems (ASes) to destination IP addresses.



• Only carry traffic if it earns you money.



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Implications & Deployment Challenges



Traffic Attraction Attacks on:



Defensive Filtering

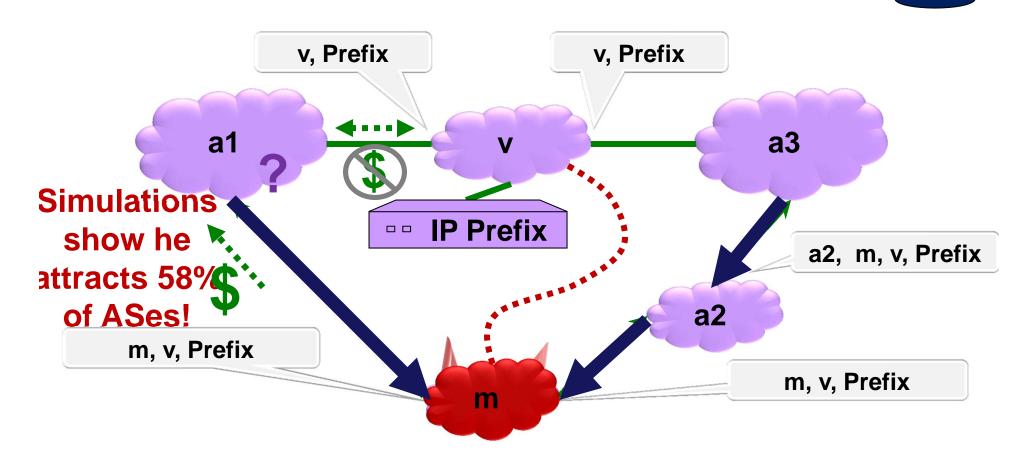
Attacker wants max number of ASes to route thru its network. (For eavesdropping, dropping, tampering, ...) v, Prefix v, Prefix victim **a**3 **a**1 **IP Prefix** a2, m, Prefix Simulations a2 \odot show he m, Prefix m, Prefix % m of ASes!

A model of routing decisions:

- Prefer cheaper paths. Then, prefer shorter paths.
- Only transit traffic if it earns you money, ie. for customers.

Proposed Security Mechanism: Origin Authentication

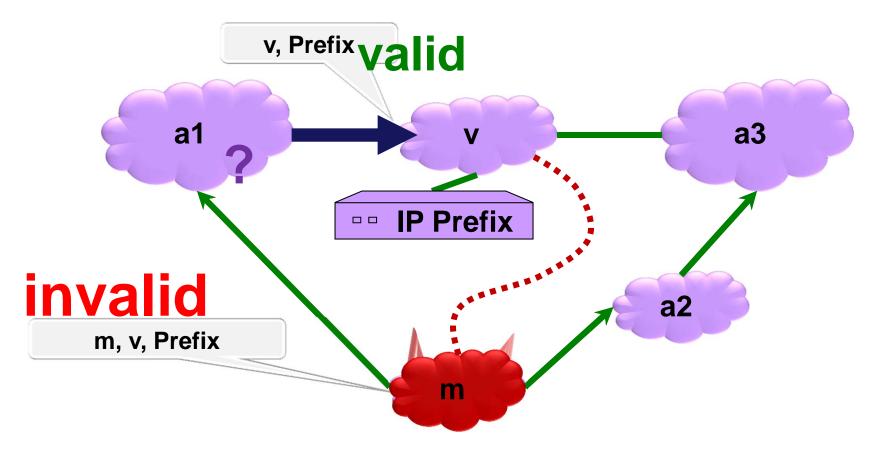
RPKI: A secure database that maps IP Prefixes to owner ASes.



Smart Attack Strategy: Announce the shortest path I can get away with to all my neighbors.

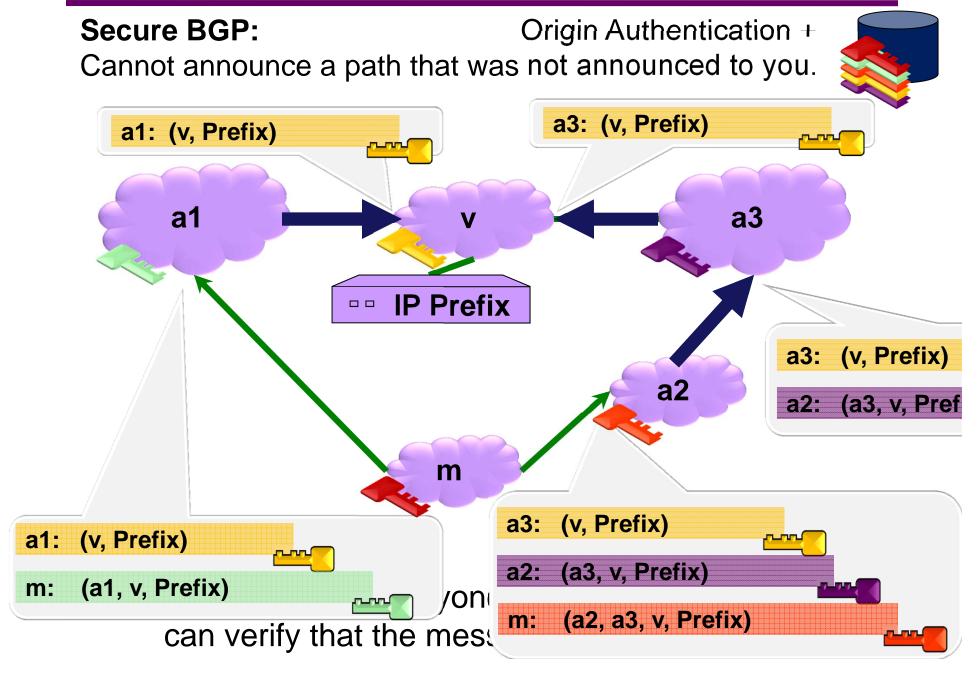
Proposed Security Mechanism: secure origin BGP

RPKI: A secure database that maps IP Prefixes to owner ASes. **soBGP**: A database of all the links in the AS-level topology.



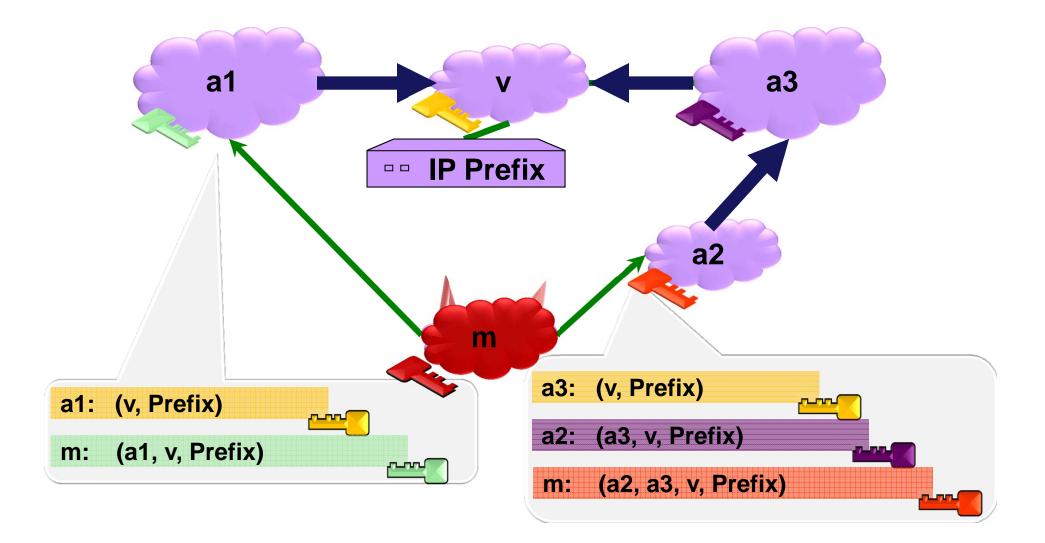
Smart Attack Strategy: Announce the shortest path I can get away with to all my neighbors.

Proposed Security Mechanism: "Secure BGP" [KLS98]



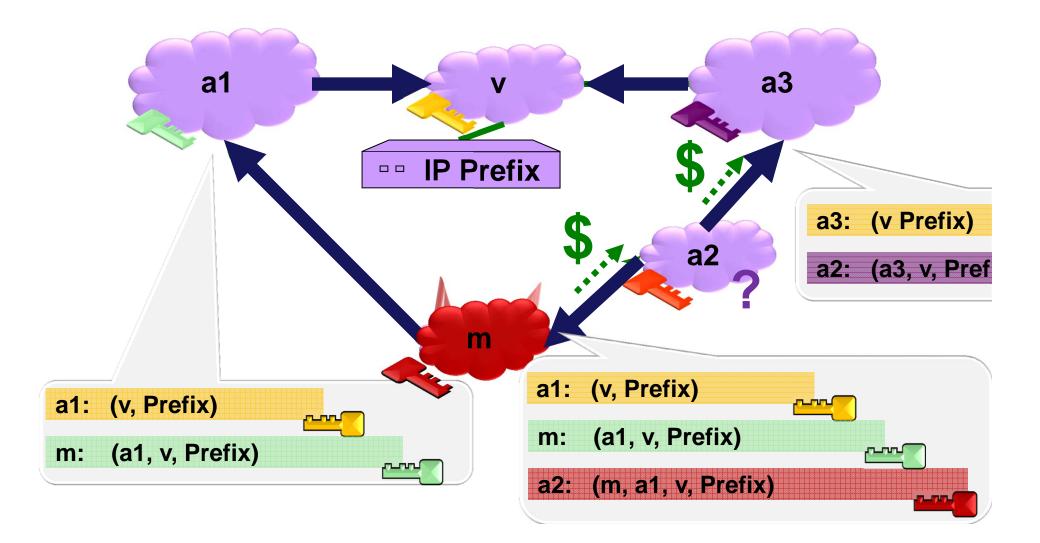
Are attacks still possible with Secure BGP? (1)

Smart Attack Strategy: Announce the shortest path I can get away with to all my neighbors!



Are attacks still possible with Secure BGP? (2)

Smart Attack Strategy: Announce the shortest path I can get away with to all my neighbors!

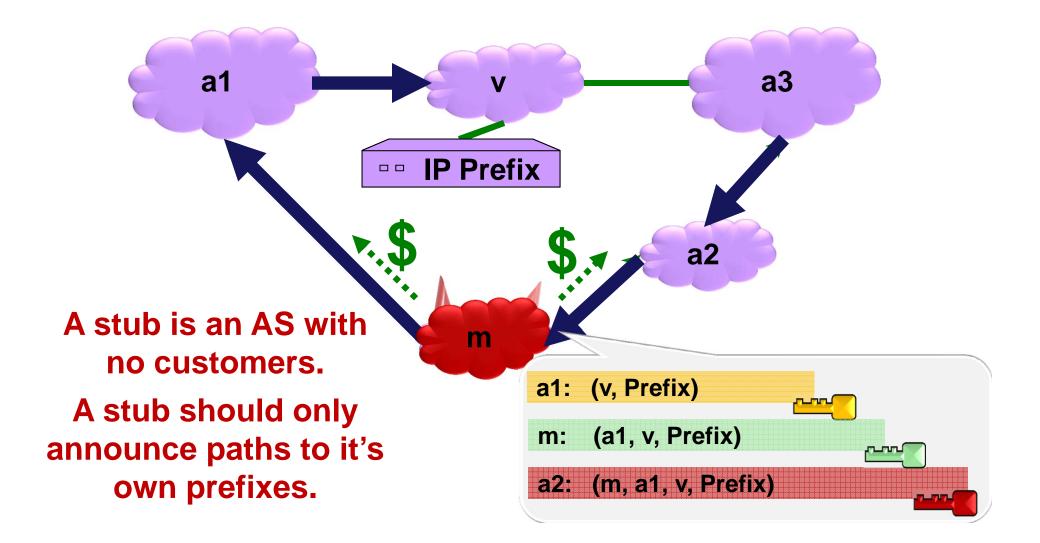


Are attacks still possible with Secure BGP? (3)

Smart Attack Strategy: Announce the shortest path I can get away with to all my neighbors! a3: (v, Prefix) a3 **a1 IP** Prefix a2 \bigcirc (v, Prefix) a1: **Simulations** (a2, v, Prefi m: m show he (m, a1, v, Pr a2: attracts 16% of ASes! (a2, m, a1, v a3:

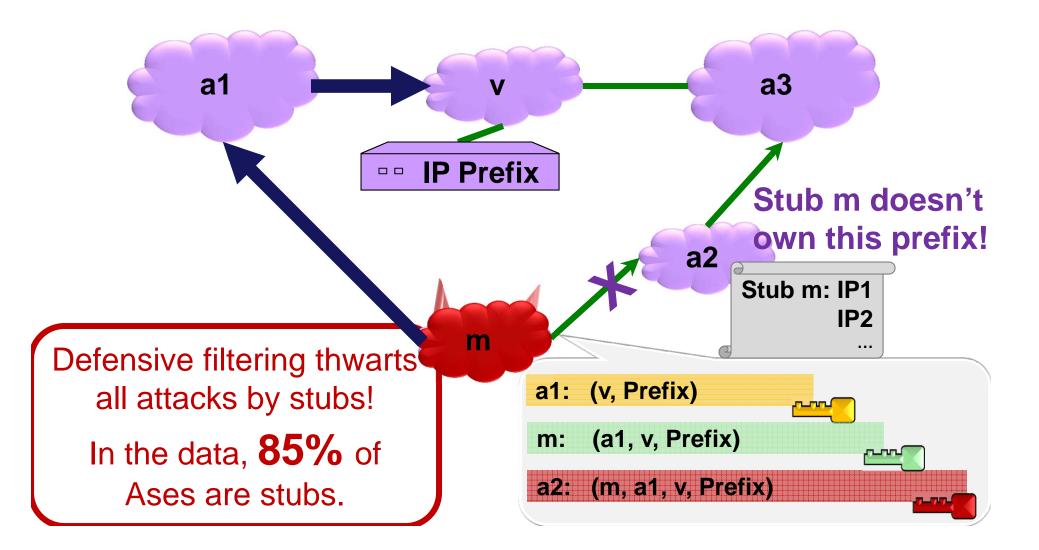
Wait! Why is this an "attack"?

Smart Attack Strategy: Announce the shortest path I can get away with to all my neighbors!



Security Mechanism: Defensive Filtering (of Stubs)

Defensive Filtering: The provider drops announcements for prefixes not owned by it's **stubs**.





This talk

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How Internet Routing Works (and why economics matter)



Secure Routing Protocols and Attacks

Theory Interlude

Results!



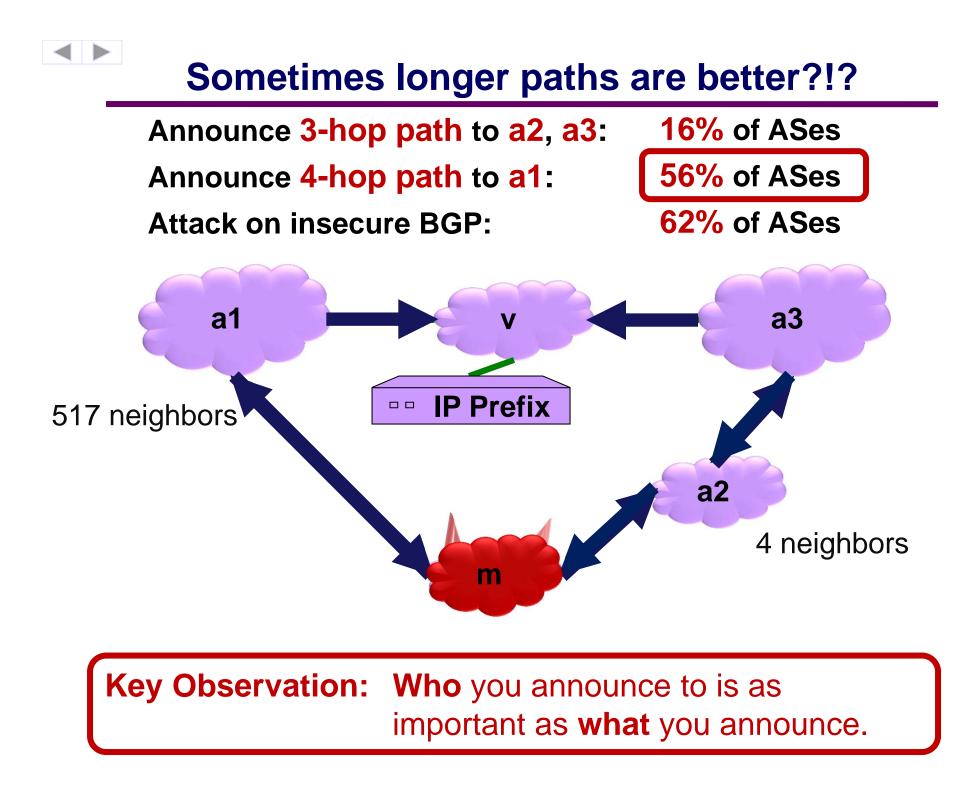
Implications & Deployment Challenges

Wait! Is this the "best" attack strategy?!?

Can't lie about my business relationship with a2, so I might as well **a**2 announce the shortest path I can. m But Not Optimal ! Smart Attack Strategy: Announce the shortest path I can get away with to all my neighbors! Sometimes Sometimes announcing to longer paths fewer neighbors is better! are better!

Theorem: It's NP hard to find the optimal attack strategy.

→ Smart Attack Strategy **underestimates** damage.





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Secure Routing Protocols and Attacks

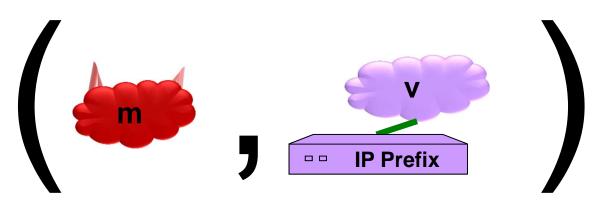
Theory Interlude

Results!



Implications & Deployment Challenges





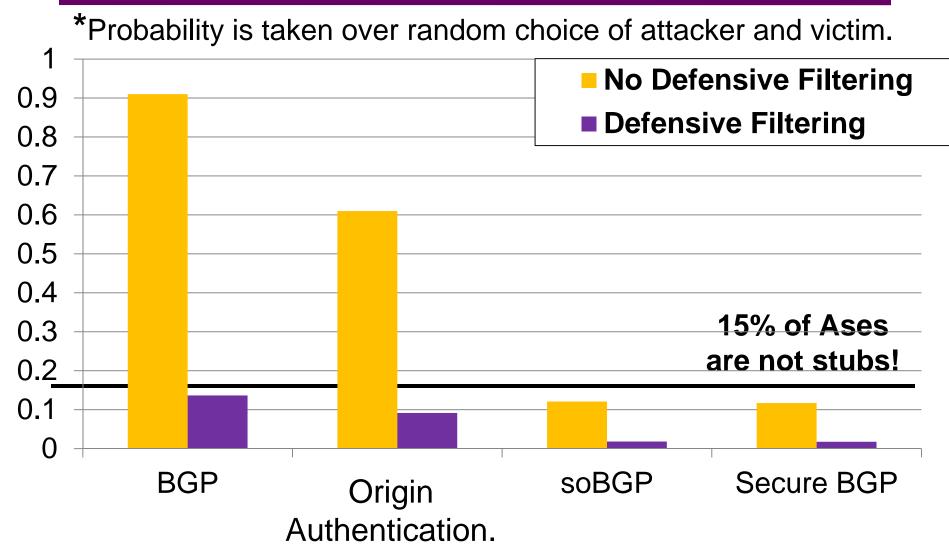
We ran multiple experiments

- For each, randomly chose (attacker, victim) pair, and
- ... simulate **Smart Attack** on each security protocol.

In the following graph:

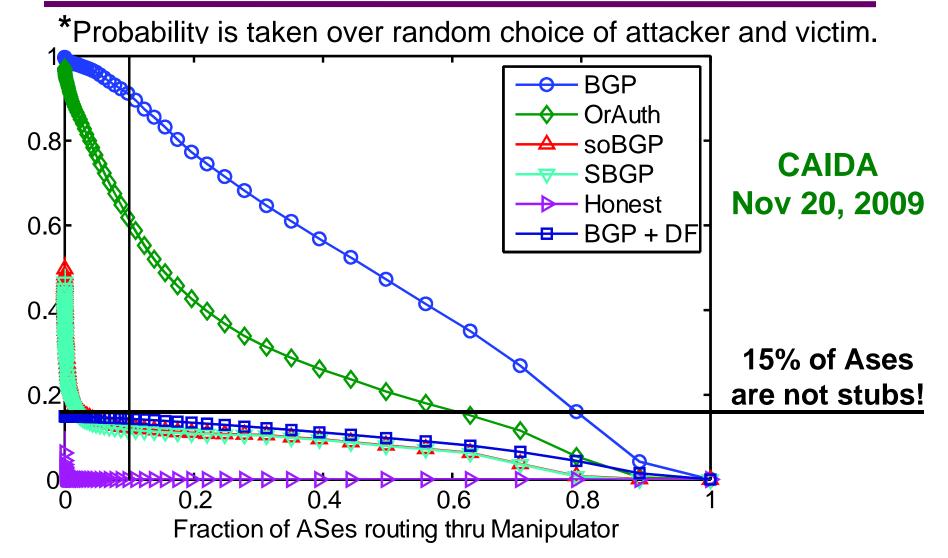
- An attacker is "successful" if it attracts **10%** of ASes.
- What fraction of pairs have a successful attacker?

Probability* Smart Attack attracts 10% of ASes



Recall that the Smart Attack Strategy underestimates damage.

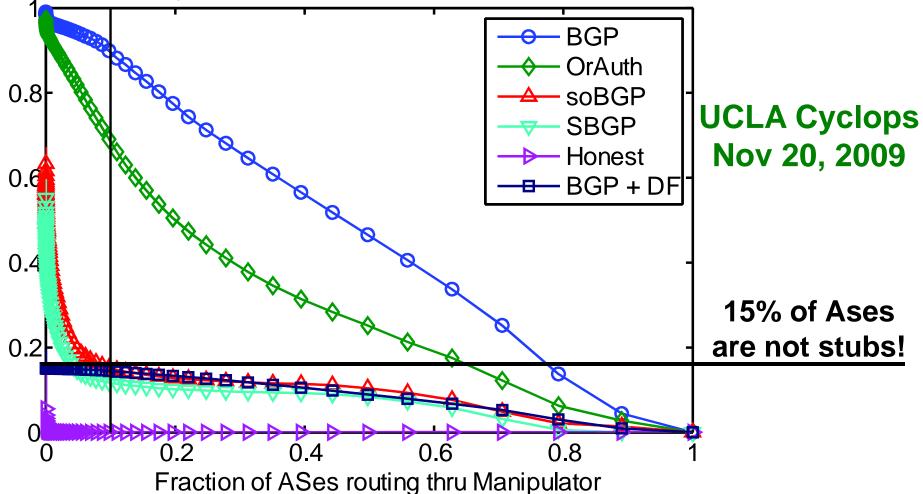
Probability* Smart Attack attracts >x% of ASes (1)



Recall that the Smart Attack Strategy underestimates damage.

Probability* Smart Attack attracts >x% of ASes (2)

*Probability is taken over random choice of attacker and victim.

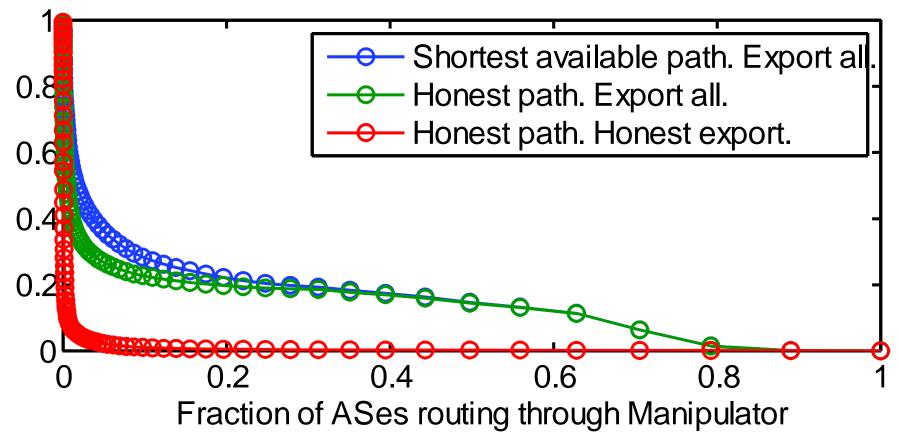


Recall that the Smart Attack Strategy underestimates damage.

The Importance of Aggressive Export Policies

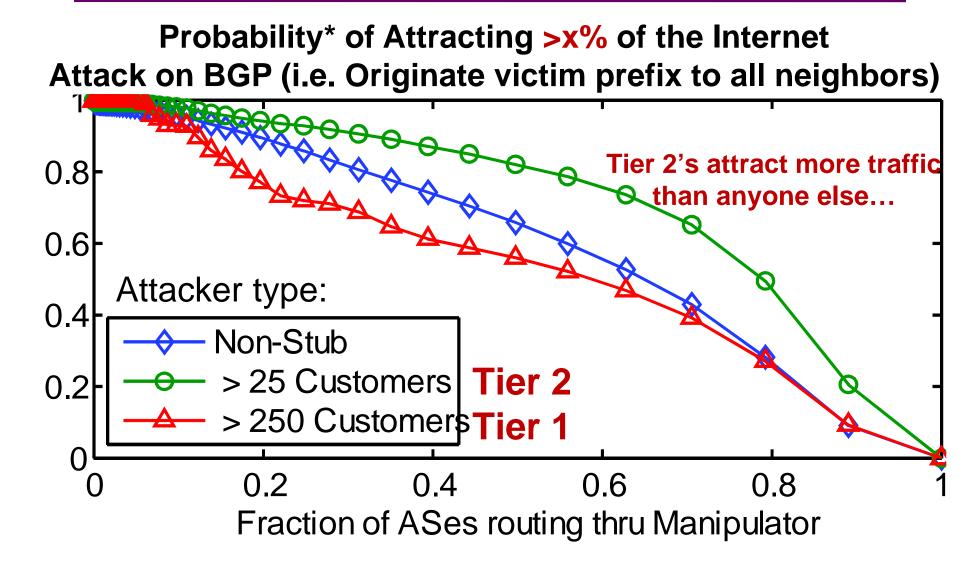
Probability* of Attracting >x% of the Internet

*Probability is over random victim and attacker with > 25 customers.



Key Observation: Who you announce to is as important as what you announce.

Tier 2's are the most effective attackers



*Probability is over random victim and attacker from different classes



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Implications & Deployment Challenges



WHO you announce to is as important as WHAT you announce

Defensive filtering is as effective as Secure BGP.

- Each mitigates a different attack strategy
- Secure BGP limits path-shortening attacks
- Filtering prevent stubs from announcing paths too widely

Why is it so hard to implement these things in practice?



Implementing Defensive Filtering ?

Today: The provider locally keeps a list of the prefixes that its stubs own.

Issues:

- 1) Relies on altruism & trust.
- 2) Maintaining prefix lists is hard.

But, some good news:



ASes.

20984

Verizon

IP2

My stub doesn't

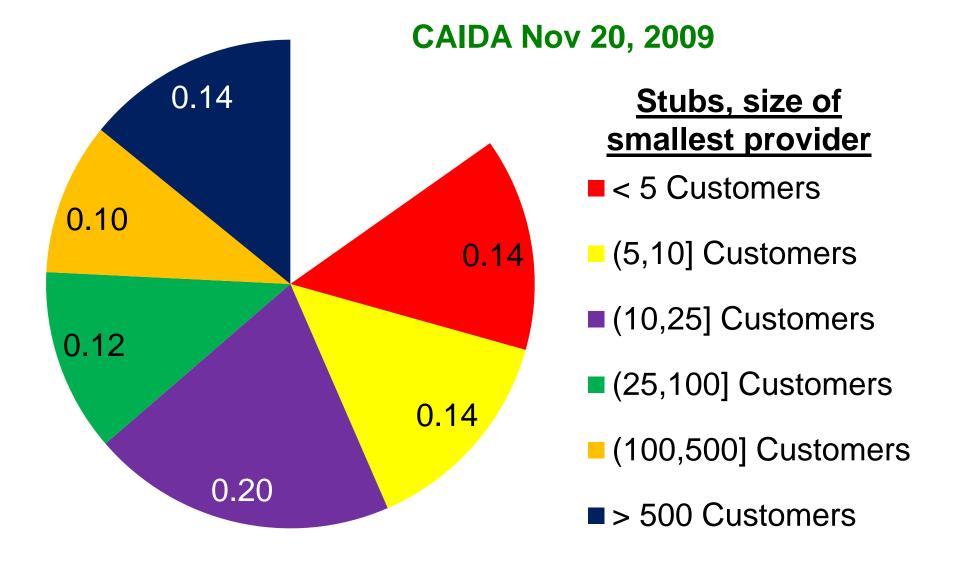
own this IP prefix!

Stub 20984: IP1

Being deployed as RPKI!

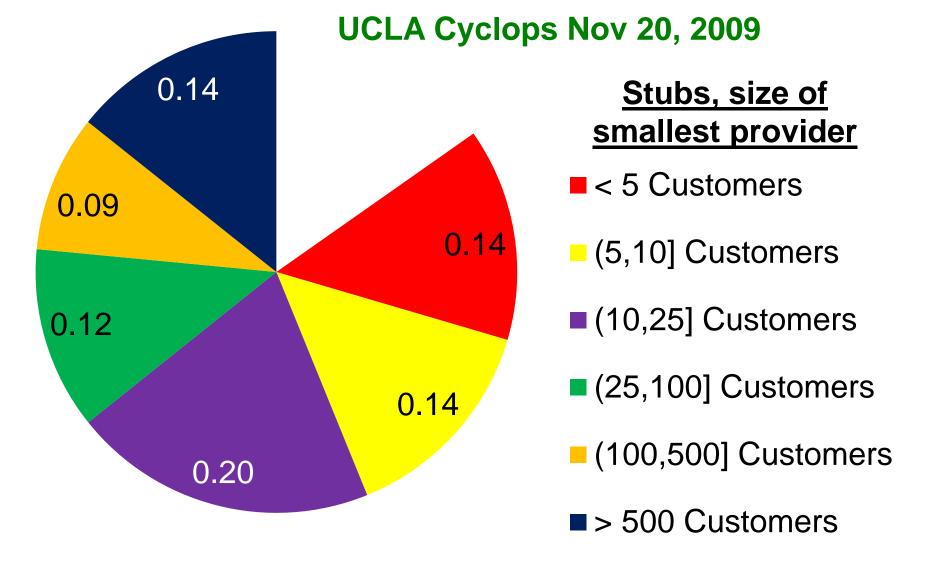
(For past few months?) prefix lists can be derived from RPKI!

What if only large ASes implement prefix lists? (1)



If ISPs with > 10 customers filter, 56% of attacks stopped.

What if only large ASes implement prefix lists? (2)



If ISPs with > 10 customers filter, 55% of attacks stopped.